

CSC/MAT 335 - EXAM 1 PRACTICE-TEST
Dr. A.J. Pounds
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Name _____

Section _____

Show all work and draw a box around numerical answers.

1. For the following calculation:

$$\frac{(12.65)(2.588)}{(13.01)} - 2.516$$

- (a) compute the result using four digit rounding.

- (b) compute the result using four digit chopping.

- (c) is there any loss of precision in the final result? If so, how many digits are lost?
If not, why not?

2. The Maclaurin series expansion of $\ln(x + 1)$ truncated to second order with its associated error term is given in Eqn. 1.

$$M_2(x) = \underbrace{x - \frac{x^2}{2}}_{\text{expansion}} + \underbrace{\frac{x^3}{3(\xi(x) + 1)^3}}_{\text{error term}} \quad (1)$$

(a) What is the value of $M_2(.25)$?

(b) What is the error bound in $M_2(.25)$?

(c) What is the absolute error in $M_2(.25)$ compared to $\ln(0.25 + 1)$?

3. You want to determine the **negative** root for the function

$$f(x) = x^3 + 2x + 10$$

using the bisection method.

- (a) Identify two values that you know bound the root.

- (b) Using the bounds from part (a), carry out four iterations of the bisection algorithm to determine a reduced domain in which the root exists.

- (c) Restarting from one of your bounds in part (a), carry out four iterations of Newton's method to find an approximation to the root.

(d) For your equation

$$f(x) = x^3 + 2x + 10$$

- i. recast the equation in the form $g(x)$ so it can be solved by the fixed point (or Picard's) method.

- ii. write a function code segment (or a subroutine) that solves for the root using the fixed point (or Picard's) method to 10 times *IEEE* single precision (32 bit) machine tolerance.

4. *NVIDIA* introduced a half precision data type (16 bits total with 1 sign bit, 5 bits in the exponent, and 10 bits in the mantissa). This later became the IEEE 754 half-precision binary floating point format known as **binary16**. What is the machine number for π in this format? Please note that the most significant and least significant bits are noted with an “m” and “l” respectively.

$$\begin{array}{c|c|c} 1 & 10001 & 1000000001 \\ \hline & \text{m} \quad \text{l} & \text{m} \qquad \qquad \text{l} \end{array}$$